SYNACKTIV

Attacking Safari in 2022

Who am I?

Quentin Meffre (@0xdagger)

- Security researcher at Synacktiv
- Vulnerability research & Exploitation

Synacktiv

- Offensive security company
- +100 ninjas
- We are hiring!



Introduction

Full chain on iPhone using the browser as entry point





Introduction

Steps to compromise Safari on the iPhone

- addrOf/fakeObj
- Arbitrary R/W
- Bypass PAC/APRR
- Overwrite JIT page code
- Arbitrary code execution!
- Apple hardened each step of a Safari exploit...









SEPARATED_WX_HEAP

The JIT page is mapped twice

- One has protections RX
- Second has protections RW

A function is jitted to copy data in the JIT page

- The function is on a page with X only protection
- The address of the RW JIT page is inlined in this function





SEPARATED_WX_HEAP

Public bypass still works with this mitigation¹

- Build an arbitrary call primitive
 - ROP/JOP
- Call the jitWriteSeparateHeaps function
- Write arbitrary code in the JIT page
- Profit!

1: https://www.sstic.org/media/SSTIC2019/SSTIC-actes/WEN_ETA_JB/SSTIC2019-Article-WEN_ETA_JB-benoist-vanderbeken_perigaud.pdf

APRR

- Hardware mitigation
- SEPARATED_WX_HEAP is replaced by APRR on supported hardware
- Atomically switches the JIT page protections using a System Register

RX ightarrow RW ightarrow RX	
MOVK	X0, #0x C110
MOVK	X0, #0xFFFF,LSL#16
MOVK	X0, #0xF,LSL#32
MOVK	XO, #0,LSL#48
LDR	X0, [X0]
MSR	#6, c15, c1, #5, X0
ISB	

APRR

Hard jump in the middle of the function¹

- The System Register value comes from a R only page shared with the kernel
- The system register value and the value from the **R only** page are compared
 - Difference \rightarrow crash

Without CFI can be bypassed like SEPARATED_WX_HEAP

1: https://github.com/phoenhex/files/blob/master/exploits/ios-11.3.1/pwn_i8.js

MOVK	X0, #0xC110
MOVK	XO, #0xFFFF,LSL#16
MOVK	X0, #0xF,LSL#32
MOVK	X0, #0,LSL#48
LDR	X 0, [X 0]
MSR	#6, c15, c1, #5, X0
ISB	
MOVK	X1 , #0xC110
MOVK	X1, #0xFFFF,LSL#16
MOVK	X1, #0xF,LSL#32
MOVK	X1, #0,LSL#48
LDR	X8, [X1]
MRS	X10, #6, c15, c1, #5
CMP	X8, X10
B.NE	loc_18BA4E060

GigaCage

TypedArray are JavaScript objects

- Often used to build arbitrary R/W
- TypedArray are allocated in a 32GB zone
 - Followed by another 32G zone allocated with **PROT_NONE**

The data buffer is now an offset to the cage and no more an address



Cannot R/W outside of the cage anymore...

GigaCage bypass

Many public documentation about the GigaCage¹

- Some public bypasses still work...
- One known bypass is to use other objects
 - More on this later in this presentation
- GigaCage is not enabled anymore on latest iOS versions
 - But attackers still can't use TypedArray to build arbitrary R/W...



StructureID randomization



JSObject inherits from the JSCell object



SYNACKTIV

The StructureID is an index

- Used to get the *Structure* of a JSObject
- Invalid StructureID → crash
- Before randomization the StructureID was incremental
 - Easy to guess a valid StructureID
 - Build fake objects without crashing

After StructureID randomization



Randomization is added to the StructureID

1 Nuke Bit | 26 StructureIDTable index bits | 5 entropy bits |

Signature is checked every time a JSObject property is accessed...

- ... but sometimes it is not!¹
- Leads to StructureID randomization bypass

StructureID randomization has been removed

StructureID uses low 32 bits of Structure address

1: https://i.blackhat.com/eu-19/Thursday/eu-19-Wang-Thinking-Outside-The-JIT-Compiler-Understanding-And-Bypassing-StructureID-Randomization-With-Generic-And-Old-School-Methods.pdf

- Pointer Authentication Code
- Hardware mitigation
- Introduced in ARMv8.3-A
- Prevents an attacker from corrupting sensitive pointers
 - Signature is added to some pointers
 - Corrupting a pointer without signing it correctly often leads to a crash

New ARM instructions used in Safari

- PAC*: Add signature to a pointer
- AUT*: Check and remove signature from a pointer
- **XPAC***: Remove signature from a pointer
- RETA*: Check X30 with context SP and return to X30 if the signature is correct
- BRA* / BLRA*: Check signature and branch



Two kinds of pointers can be signed

- Data
- Instruction

Two keys can be used for each kind

- Key A
- Key B

A context is often used to avoid pointer substitution

• A pointer can also be signed with a null context...



The signature is stored in the top bits of a pointer

The signature length depends on the key/pointer kind

- 16 bits
- 24 bits





Instruction pointers

- VTable function pointer => PACIA
- Return value stored on the stack => PACIB
- JIT Code pointer => PACIB

Data pointers

- VTable pointer => PACDA
- Sensitive data pointer (TypedArray data pointer...) => PACDB
- JIT instructions => PACDB



What is not signed in Safari?



PAC bypass

Bypassing PAC is a security issue in itself

• Apple takes **PAC** bypasses very seriously

Many PAC bypasses have been disclosed since PAC introduction

- Apple fixes each of them
 - Hardware improvement
 - Software improvement



PAC bypass: design issue

If a pointer authentication fails

- Signature is removed and one of the top bits is flipped
- Does not raise an exception

If the pointer is signed again after the failed AUT*

- Correct signature is added, with a flipped bit
- PAC bypass: flip the bit again to get the correct signature
- EnhancedPAC is implemented first on A14 SoC
 - Signing invalid pointers will discard the signature
 - Can't leak the signature anymore...

PAC bypass: bruteforce

- The signature can still be bruteforced...
- ...but Apple killed this bypass again
- The compiler option -fptrauth-auth-traps is used
 - Adds a check after all AUT* instructions
 - If the signature given to the AUT^* instruction is invalid $\rightarrow ABORT$

AUTIB	X16, X17
MOV	X17, X16
XPACI	X17
CMP	X16, X17
B.EQ	loc 18BA4ABD8
BRK	#0xC471
SYNACKTIV	

PAC bypass: bruteforce

Apple added a new feature in the A15 SoC

- ARMv8.6-A FPAC extension
- If an AUT* instruction fails, an exception is now raised

Apple killed this exploitation method with this feature



PAC bypass: null context chained



- Initially, many pointers were signed with a null context
- A potential bypass could be to use null signed pointers in a JOP chain

- Build powerful primitives
- Never seen publicly
- Since iOS 15 this attack has been almost killed
 - Very few pointers are still signed with a null context



PAC bypass

More bypasses¹

- Unprotected code pointers
- Race condition with the JIT thread
- Blocking the JIT thread while copying data on the JIT page
- Signal handlers corruption
- All of these bypasses have been fixed

1: https://googleprojectzero.blogspot.com/2020/09/jitsploitation-three.html



PAC R/W

- PAC doesn't sign a lot of sensitive data pointers
- Some object can be wrapped into a JSObject
 - DOMRect
 - Contains 4 doubles
 - Has methods to read and write these doubles
- Faking a wrapper to a DOMRect object
 - Arbitrary R/W

Method used by a public exploit¹

1: https://blog.google/threat-analysis-group/analyzing-watering-hole-campaign-using-macos-exploits/



PAC kill R/W method

Method killed by iOS 15.4

- Some wrappers to sensitive wrapped objects are now signed
 - Most of them manipulate floats/doubles
 - Killed many arbitrary R/W methods

Introduce SignedPtrTraits which enables Ref pointers to be protected with PtrTags.

JIT Code signature



The JIT compilation can be done in another thread

- The assembly code is stored in a temporary buffer while doing compilation
- The temporary buffer content is copied in the JIT page at the end of the compilation

Before JIT code signature

- Race the JIT thread to put arbitrary code in the temporary buffer
- Profit!
- But...



JIT Code signature

Apple introduced the JIT code signature

- Stop attackers from overwriting the *JIT* code buffer
- Software mitigation based on PAC
- Instructions stored in the temporary buffer are signed
 - Each instruction signature generates a hash stored in the hash buffer
 - Signed with previous hash and **PACDB**
- Signature is checked when the temporary buffer is copied in the JIT page
 - If the signature is invalid \rightarrow Crash



JIT Code signature PIN

- The hash used to sign the next instruction was not protected
- It is now signed with a unique identifier (PIN)
 - Each JIT compilation uses a JIT Page different PIN
 - PIN informations are stored in the JIT page
 - An attacker can't modify them





The A15 SoC brings a new complex mitigation

• The **JITCage**!

- The JITCage stops attackers from calling arbitrary functions from the JIT page
- The JIT page is now mapped with a new flag

• MAP_JITCAGE?

The XNU open-source project doesn't have references about this flag...

...but the KernelCache has references!

```
void fastcall enable jitbox ( int64 thread)
 int64 current thread; // x0
 int64 v3; // x8
 unsigned __int64 StatusReg; // x9
 current_thread = get_thread_ro();
 if ( current thread == thread )
   v3 = *( OWORD *) (current thread + 0x358);
   StatusReg = _ReadStatusReg(ARM64_SYSREG(3, 0, 13, 0, 4));
   *(_QWORD *)(*(_QWORD *)(StatusReg + 0x158) + 0x218LL) = v3;
   *(_QWORD *)(*(_QWORD *)(StatusReg + 0x158) + 0x210LL) = *(_QWORD *)(current thread + 0x350);
   WriteStatusReg(ARM64 SYSREG(3, 4, 15, 15, 4), *( QWORD *) (current thread + 0x358));
   WriteStatusReg(ARM64 SYSREG(3, 4, 15, 15, 1), *( QWORD *)(current thread + 0x350));
   isb(0xFu);
```


The kernel sets new System Registers using

- The size of the JIT page
- The address of the JIT page
- Some unknown flags
- The KernelCache has no other information
- The interesting part of the JITCage is implemented in the A15 SoC

The following instructions can't be executed in the JITCage

- RET
- BR/BLR/BL
- SVC
- MRS/MSR

If one tries to execute these instructions in the JITCage

The processor raises an EXC_BAD_INSTRUCTION exception



- The PAC IA/IB keys are different in the JITCage
- Can't sign instruction pointers in the JITCage
 - **PACIA** doesn't add signature if executed in the JITCage
 - PACIB can only sign pointer that points into the JITCage
 - PACD* seems unaffected by the JITCage

- The JIT code has to call functions outside of the JITCage
- Setting a System Register allows changing IA key
 - Instruction pointers used by the JITCage are signed with the IA key
- Only done once when the JavaScript engine is initialized
- Can't be done anymore after
 - MRSX8, #4, c15, c15, #6ORRX8, X8, #0x8000MSR#4, c15, c15, #6, X8
- An attacker can't easily call functions outside of the JITCage

Conclusion 1/2

Getting arbitrary code on latest iPhone involves finding:

- A vulnerability
- A new method to build arbitrary R/W
- A PAC bypass
- An **APRR** bypass
- A JITCage bypass

One solution for attackers could be to implement the next stage using JavaScript only...

Conclusion 2/2

2022 in short

- Yet another mitigation
- Yet other exploitation methods killed

What to expect in the next years?

Same as above?

Maybe it's time for attackers to find another entry point than the browser...

- ...or maybe not? :-)
- JavaScript is a powerful engine to attack all those mitigations

